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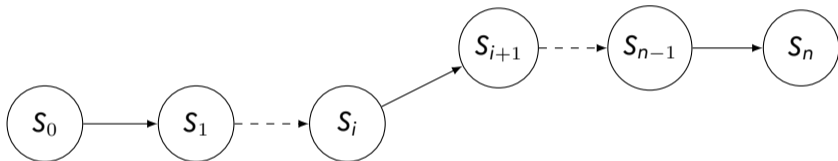
Security of Complex Software

How to Move from Characterising the Fault Effects to Exploitation?

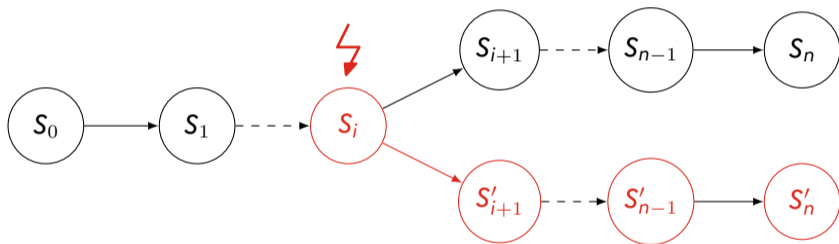
Guillaume Bouffard

Co-joint work with: Alexandre Iooss and Thomas Troughkine
Agence nationale de la sécurité des systèmes d'information

Fault Injection Attacks: Effects on Software



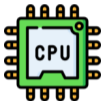
Fault Injection Attacks: Effects on Software



Scenario of a Fault Injection Attack



Target component



Fault
models



Binary/Source-code



Attack
paths



Exploitation

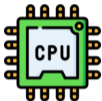


Fault injection medium

Scenario of a Fault Injection Attack



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Exploitation

Characterisation

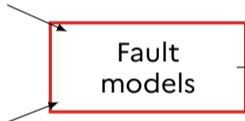
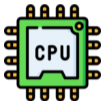
Analysis

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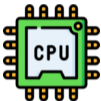
Exploitation

Scenario of a Fault Injection Attack

Characterisation approaches:

- From software POV: T. Trouchkine's PhD Thesis [Tro21].
- From hardware logic: A. Marotta's PhD Thesis [Mar].

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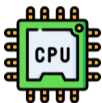
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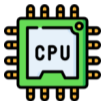
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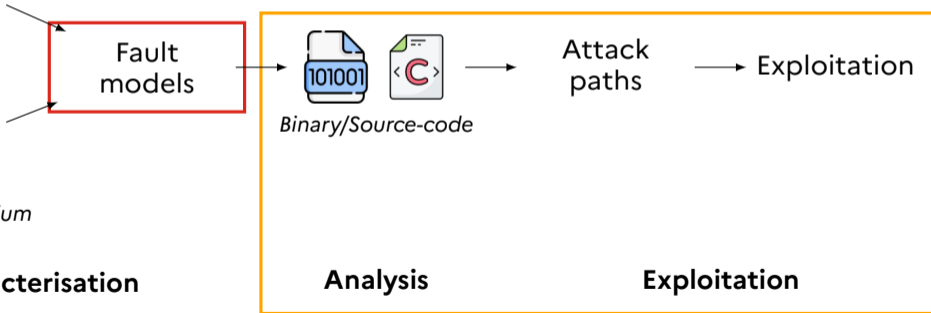
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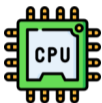


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Security analysis of sudo [Tro21, ITB23].

Fault Injection Attacks



To succeed a Fault Injection (FI) attack, one needs to:

- Know **where** attacks? \Rightarrow FI medium (x, y, z, θ) position on the target component;
- Know **how** attacks? \Rightarrow FI medium parameters;
- Know **when** attacks? \Rightarrow Time t to confuse target program.

Fault Injection Attacks (cont.)



- 1 **Where?/How?:** Characterize the sensitivity of the target component to a fault medium [TBC21, PHB⁺19]
- 2 **When?:** Transfer this sensitivity to a target application [GHHR23, Dur16]
- 3 **Conducting an attack through exploitation** [Wer22, TBE⁺21]

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Can one determine **when** to induce a fault in a complex software by characterizing (**where/how**) the effect of a fault?



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Can one determine **when** to induce a fault in a complex software by characterizing (**where/how**) the effect of a fault?

Target: The sudo application on a Raspberry Pi 4 running the Raspberry Pi OS¹ (derived from Debian).

¹See: <https://www.raspberrypi.com/software/>

1. Characterisation

Analysis of fault effects on a complex CPU [Tro21]



During a fault, at least one microarchitecture block is disturbed.

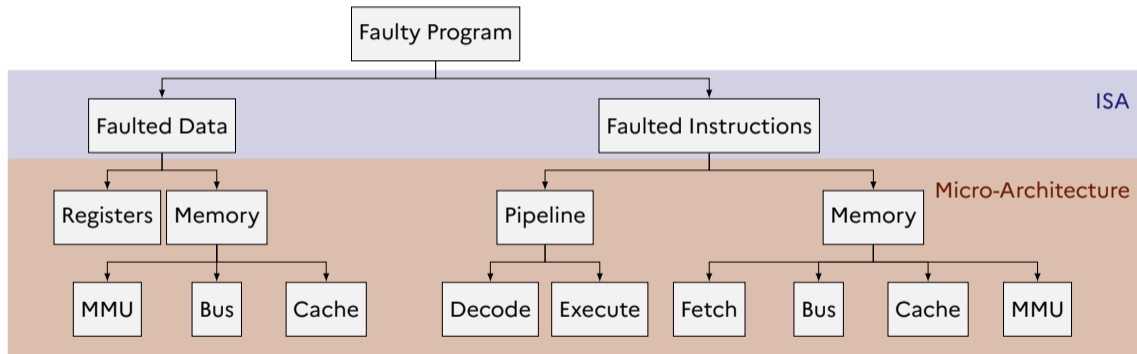


During a fault, at least one microarchitecture block is disturbed.

Modus Operandi

- A test program is faulted during its execution.
- Fault different test programs to gather information on the behavior of microarchitectural blocks.

Analysis of fault effects on a complex CPU [Tro21]



Test program 1

```
orr r4, r4;
/*
 * Arbitrary number
 * of repetitions
 */
```

```
orr r4, r4;
```

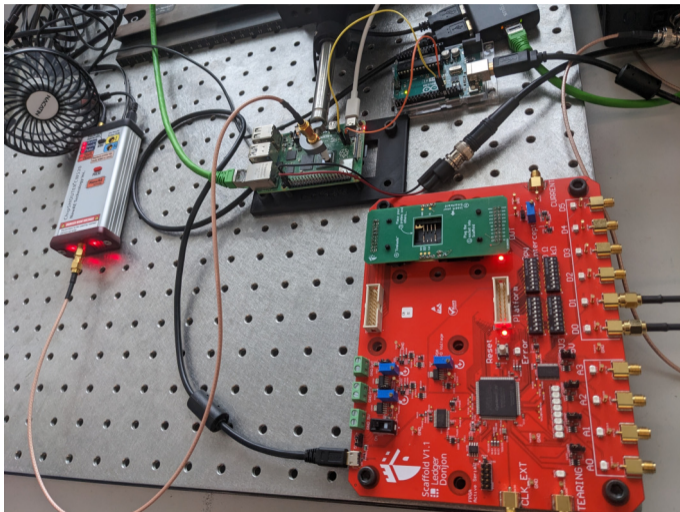
Test program 2

```
and r4, r4, #255;
/*
 * Arbitrary number
 * of repetitions
 */
and r4, r4, #255;
```

Table: Initial values of the registers.

Register	Initial value for ORR R4, R4 repetition	Initial value for AND R4, #255 repetition
r0	0xFFFE0001	0xFFFE0001
r1	0xFFFD0002	0xFFFD0002
r2	0xFFFB0004	0xFFFB0004
r3	0xFFF70008	0x000000FF
r4	0xFFEF0010	0xFFEF0010
r5	0xFFDF0020	0xFFDF0020
r6	0xFFBF0040	0xFFBF0040
r7	0xFF7F0080	0xFF7F0080
r8	0xFEFF0100	0xFEFF0100
r9	0xFDF0200	0xFDF0200

Test Bench



Sensitivity Map and Results

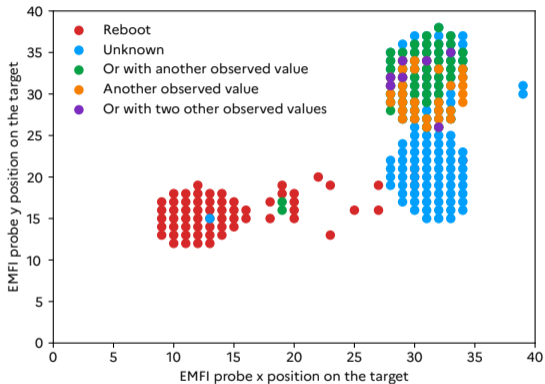


Figure: ORR R4, R4

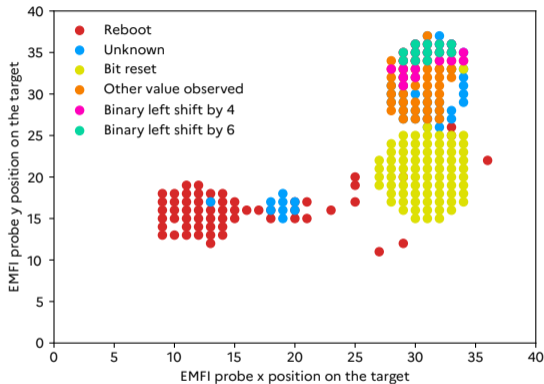


Figure: AND R4, R4, #255

Sensitivity Map and Results

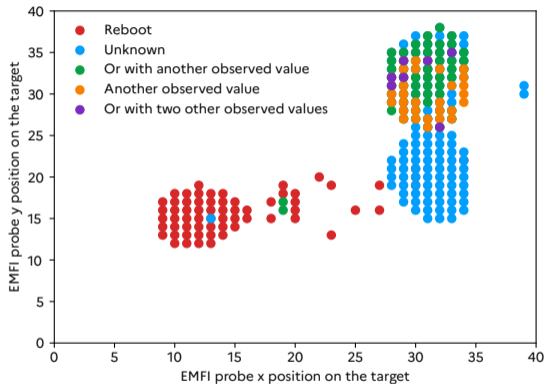


Figure: ORR R4, R4

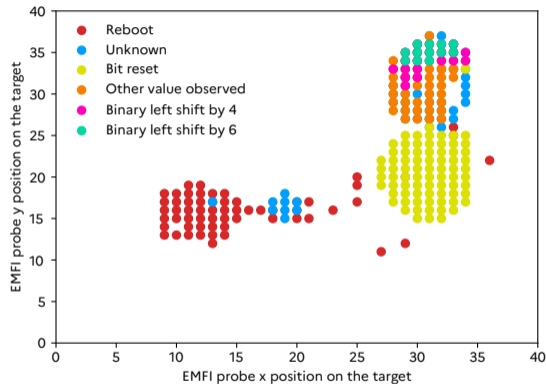


Figure: AND R4, R4, #255

👉 Setting the 8 least significant bits to 0 in the executed instructions.

2. Fault Model Transposition



Simulation Goal: Determine **when** to inject the fault into the execution flow to achieve a desired effect. We focus on simple fault.

Target: Binary sudo from Raspberry Pi OS² (based on Debian)

²https://downloads.raspberrypi.org/raspios_lite_armhf/images/raspios_lite_armhf-2022-09-26/



Simulation Goal: Determine **when** to inject the fault into the execution flow to achieve a desired effect. We focus on simple fault.

Target: Binary sudo from Raspberry Pi OS² (based on Debian)

Authentication rules in /etc/pam.d/common-auth

```
# here are the per-package modules (the "Primary" block)
auth      [success=1 default=ignore]      pam_unix.so nullok
# here's the fallback if no module succeeds
auth      requisite                        pam_deny.so
```

sudo calls PAM, which uses the pam_unix.so module.

²https://downloads.raspberrypi.org/raspios_lite_armhf/images/raspios_lite_armhf-2022-09-26/

Dynamic Instrumentation of `pam_unix.so` with Rainbow



Rainbow is a fault injection simulator developed by the Ledger Donjon, based on Unicorn-Engine (QEMU).

⚠️ `sudo` is **dynamically** linked with `glibc` and PAM.

 Ledger-Donjon/rainbow



Rainbow is a fault injection simulator developed by the Ledger Donjon, based on Unicorn-Engine (QEMU).

⚠️ sudo is **dynamically** linked with glibc and PAM.

We implemented a loader based on CLE in Rainbow.

Dynamic dependencies of pam_unix.so

```
libpam.so.0, libcrypt.so.1, libselinux.so.1, libnsl.so.2,  
libtirpc.so.3, libc.so.6, ld-linux-armhf.so.3, libaudit.so.1,  
libdl.so.2, libpcr2-8.so.0, libgssapi_krb5.so.2,  
libpthread.so.0, libcap-ng.so.0, libkrb5.so.3,  
libk5crypto.so.3, libcom_err.so.2, libkrb5support.so.0,  
libkeyutils.so.1, libresolv.so.2
```



 Ledger-Donjon/rainbow



```
1 def fault_model(emu):
2     # Get PC value
3     pc = emu["pc"]
4     # Get next instruction
5     instr = emu[pc]
6
7     # Patch and run modified instruction
8     i = int.from_bytes(instr, "little") & 0xFFFF_FF00
9     instr_patched = i.to_bytes(4, "little")
10    emu[pc] = instr_patched
11    emu.start(pc, 0, count=1)
12
13    # Restore correct instruction
14    emu[pc] = bytes(instr)
```

Exhaustive Testing in Simple Fault



```

1  emu = rainbow_arm()
2  emu.load("arm-libs/pam_unix.so")
3  emu[0xE0000000] = f"toto\x00".encode()
4  emu[0xF0000000] = f"$6$hH.15uU5laaxuXH$anemvMyc.gFyc[...]nSGEO.\x00".encode()
5  emu["r0"] = 0 # pamh (used for pam_syslog calls)
6  emu["r1"] = 0xE0000000 # const char *p
7  emu["r2"] = 0xF0000000 # char *hash
8  emu["r3"] = 1 # unsigned int nullok
9  pc_stopped = emu.start_and_fault(fault_model, i, 0x00405b40, 0, count=1000)
10 print(emu["r0"]) # if 0, then auth is successful
  
```

Variant of the Ledger blog post [IS22]. We hook malloc, calloc, and free.
 22 minutes or 3.2 seconds with SHA2-256 hooked (crypt_r).

Exhaustive Testing in Simple Fault: Results

Fault Found

```

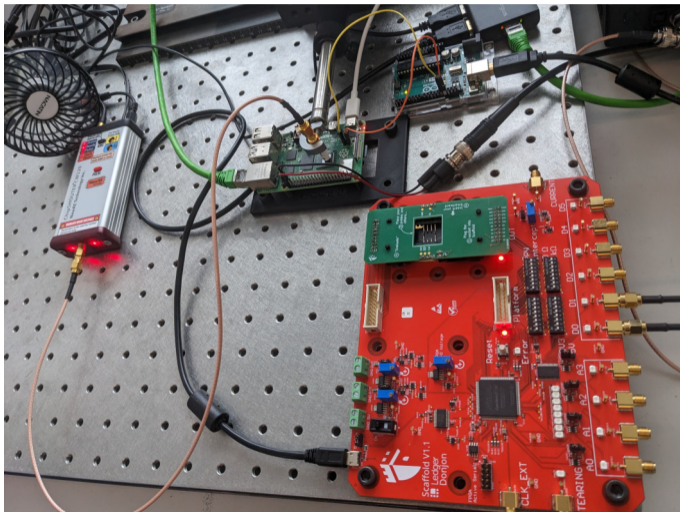
; i=1201 (in pam_unix.so)
MOVNE R4, #7
; become
MOVNE R4, #0
    
```

```

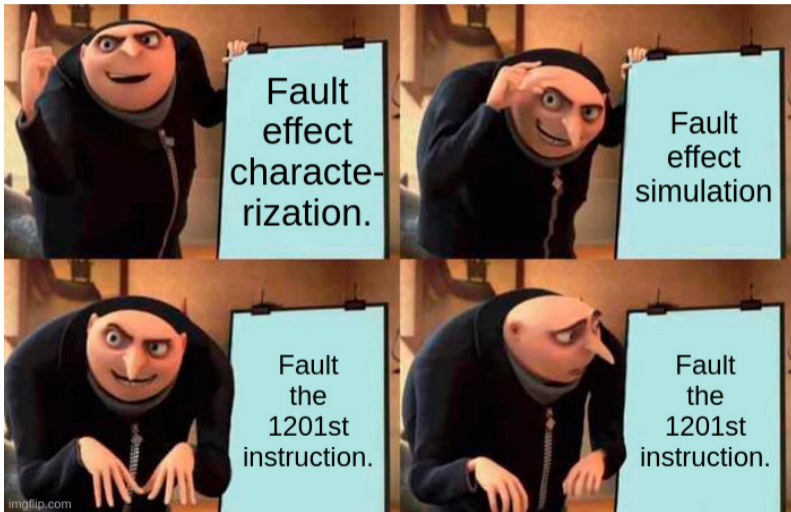
; [...]
@121: movne    r4,#0x7    ; r4 <- 0 with fault
@122: cmp      r3,#0x0
@123: cpyne   r3,r7
@124: movne   r1,#0x0
@125: beq     LAB_00405cb4
        LAB_00405ca4:
@126: strb    r1,[r3],#0x1
@127: ldrb    r2,[r3,#0x0]
@128: cmp     r2,#0x0
@129: bne    LAB_00405ca4
        LAB_00405cb4:
@130: cpy     r0,r7
@131: bl      free
@132: cpy     r0,r4    ; returns the value of r4
@133: ldmia   sp!,{r4,r5,r6,r7,r8,r9,r10,pc}
    
```

3. Exploitation

Test Bench



Transition to Reality



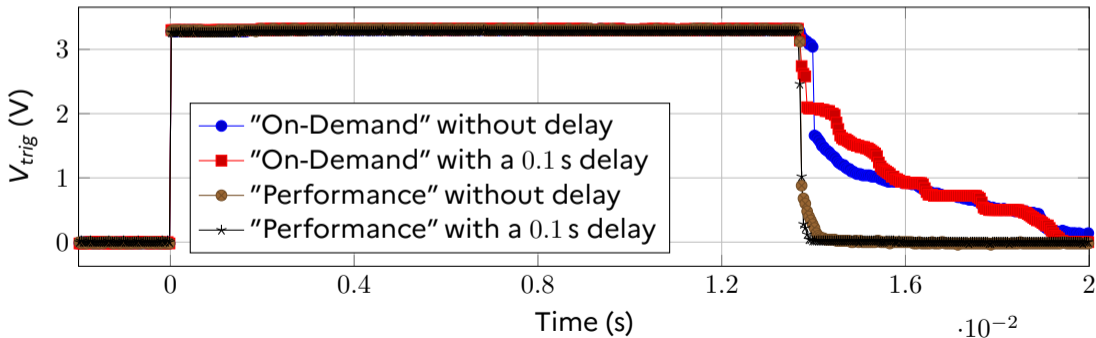


Modus Operandi

- Creation of a wrapper to launch sudo via Netcat and trigger a GPIO.
- sudo forced on CPU core 3.

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Trick to Find the Fault Instant



Use of an open-sample

```

LAB_00405c7c
00405c7c 05 10 a0 e1 cpy r1,r5
00405c80 07 00 a0 e1 cpy r0,r7
00405c84 96 ee ff eb bl <EXTERNAL>::strcmp
00405c88 00 30 d7 e5 ldrb r3,[r7,#0x0]
00405c8c 00 40 50 e2 subs r4,r0,#0x0
00405c90 02 00 a0 e3 mov r0,#0x2
LAB_00405c94
00405c94 01 20 a0 e3 mov r2,#0x1
00405c98 04 70 a0 e3 mov r7,#0x4
00405c9c 00 00 00 ef swi 0x0
00405ca0 03 00 00 0a beq LAB_00405cb4
LAB_00405ca4
00405ca4 01 10 c3 e4 strb r1,[r3],#0x1
00405ca8 00 20 d3 e5 ldrb r2,[r3,#0x0]
00405cac 00 00 52 e3 cmp r2,#0x0
00405cb0 fb ff ff 1a bne LAB_00405ca4
LAB_00405cb4
00405cb4 07 00 a0 e1 cpy r0,r7

LAB_00405c7c
00405c7c 05 10 a0 e1 cpy r1,r5
00405c80 07 00 a0 e1 cpy r0,r7
00405c84 96 ee ff eb bl libc.so.6::strcmp
00405c88 00 30 d7 e5 ldrb r3,[r7,#0x0]
00405c8c 00 40 50 e2 subs r4,r0,#0x0
00405c90 07 40 a0 13 movne r4,#0x7
LAB_00405c94
00405c94 00 00 53 e3 cmp r3,#0x0
00405c98 07 30 a0 11 cpyne r3,r7
00405c9c 00 10 a0 13 movne r1,#0x0
00405ca0 03 00 00 0a beq LAB_00405cb4
LAB_00405ca4
00405ca4 01 10 c3 e4 strb r1,[r3],#0x1
00405ca8 00 20 d3 e5 ldrb r2,[r3,#0x0]
00405cac 00 00 52 e3 cmp r2,#0x0
00405cb0 fb ff ff 1a bne LAB_00405ca4
LAB_00405cb4
00405cb4 07 00 a0 e1 cpy r0,r7
    
```

```

120 pbVar8 = 0;
121 bVar1 = pbVar8[1];
122 pbVar8 = pbVar8 + 1;
123 }
124 }
125 }
126 iVar4 = strcmp((char *)pbVar9,hash);
127 pcVar3 = (char *) (uint)pbVar9;
128 bVar10 = iVar4 == 0;
129 /* can be faulted with 0x0FFF_FF00 */
130 cVar5 = extraout_r1_00;
131 LAB_00405c94:
132 software_interrupt(0);
133 if (!bVar10) {
134 do {
135 pcVar7 = pcVar3 + 1;
136 *pcVar3 = cVar5;
137 pcVar3 = pcVar7;
138 } while (*pcVar7 != '\0');
139 }
140 free((void *)0x4);
141 return iVar4;
142 }
143
    
```

Decompile: verify_pwd... Bytes: pam_unix.o

Trick to Find the Fault Instant



Problem: The delay introduced by the `syscall` and parent process.

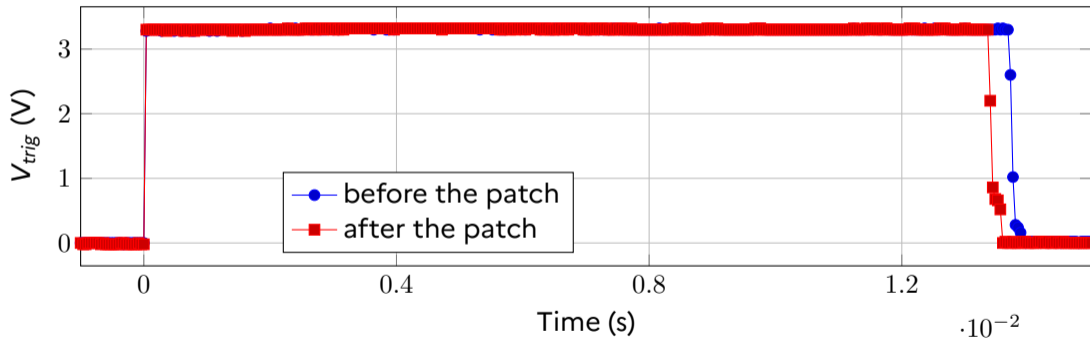


Figure: Average of 1024 calls to `sudo` with the incorrect password, CPU performance mode, delay of 0.1 s.

- The expected fault model was not consistently found at the same FI medium position.



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- Partial sudo bypass:
 - `sudo --validate`



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- Partial sudo bypass:
 - `sudo --validate` ⇒ success
 - `sudo whoami`



- The expected fault model was not consistently found at the same FI medium position.
- Partial sudo bypass:
 - `sudo --validate` ⇒ success
 - `sudo whoami` ⇒ not success



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 - `sudo --validate` ⇒ **success**
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 - The reasons for this difference between simulation and reality are currently unknown.



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- Partial sudo bypass:
 - `sudo --validate` ⇒ **success**
 - `sudo whoami` ⇒ **not success**
 - The reasons for this difference between simulation and reality are currently unknown.
- Require an open-sample to determine real-world injection timing:
 - overcoming with a cycle-accurate emulation ⇒ it's rarely available for complex systems.
 - Using side-channel simulators to identify the injection point on a simulated trace, but precise modeling of the target is currently challenging.



In this work:

- Transposition of the fault model to a complex program (sudo)
- ⚠ Transition from analyzing a binary to a bench attack
- Use of open-source tools:
 - Ledger-Donjon/rainbow (proposed patches submitted).
 - Ledger-Donjon/scaffold
 - newaetech/ChipSHOUTER



In this work:

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What's next?

- Can we do without an *open-sample*?
- How to protect against it?
- Transition to RISC-V.



Security of Complex Software

How to Move from Characterising the Fault Effects to Exploitation?

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Credit: icons downloaded from Flaticon



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- [PHB⁺19] Julien Proy, Karine Heydemann, Alexandre Berzati, Fabien Majéric, and Albert Cohen, *A first isa-level characterization of EM pulse effects on superscalar microarchitectures: A secure software perspective*, Proceedings of the 14th International Conference on Availability, Reliability and Security, ARES 2019, Canterbury, UK, August 26-29, 2019, ACM, 2019, pp. 7:1–7:10.



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